# Vertex distinguishing colorings of graphs 

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## We shall consider

edge colorings of graphs and use these colorinings to distinguish the vertices of the graphs.
I. proper colorings
(joint work with R. Kalinowski, M. Pilśniak and J.
Przybyło form AGH University)
II. general colorings
(joint work with O.Baudon and J.Bensmail from LaBRI and J. Przybyło form AGH University)

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- A graph is Class 2 if $\chi^{\prime}(G)=\Delta(G)+1$.

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## Coloring distinguishing (all) vertices

- 1. Košice (observability, obs) Master thesis by Roman Soták (1992), Černy, M.Horňák, R.Soták, Mat. Slovaca (1996) M.Horřák, R.Soták Ars Combin. (1995) M.Horňák, R.Soták Discrete Math. (1997)


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- 1+2.

Graph Theory Week, Banach Centre 1996 (Horňák + Favaron)

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- Remark. We assume that our graph has neither $K_{2}$ nor two $K_{1}$ as components.


## vdi - A conjecture

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- Example (!): complete graphs $K_{2 p}$.


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Theorem:
If $\delta(G) \geq \frac{n}{3}$, then vdi $(G) \leq \Delta(G)+5$.


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- We have $\binom{k}{2} \geq n$
- $\operatorname{vdi} \geq \sqrt{2 n}$.


## By the way ...

- Conjecture 2.
$\pi \leq \operatorname{vdi}(G) \leq \pi+1$
where $\pi=\max _{i}\left\{\min _{k}\left\{k:\binom{k}{i} \geq n_{i}\right\}\right\}$ where $n_{i}$ denotes the number of vertices of degree $i$.


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- Actually, first formulated in Soták's Master thesis (1992; in slovak).
- Still open despite of many papers: mainly by P.N.Balister with:
B.Bollobás, O.M.Riordan, R.H.Schelp, A.Kostoczka, Hao Li.

An example


## Another example


$C_{4}$


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- $W(x)$ - the set of all color walks starting from $x$,
- two vertices $x$ and $y$ are similar if $W(x)=W(y)$.
- $\mu(G)$ - the minimum number of colors in a proper edge-coloring of a graph $G$ such that no two distinct vertices are similar.


## Main result

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- R. Kalinowski, M. Piśniak, J. Przybyło and M. Woźniak, How to personalize the vertices of a graph?, European Journal of Combinatorics, 40 (2014), 116-123.


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- (main theorem for graphs class I)

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$\bigcirc \bullet \bullet \bullet \bullet \bullet \longrightarrow$



## $C_{8}$ - a coloring



## One of cases <br> 



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## One of cases

## $C_{3} \square K_{2}$


$C_{9}$

$C_{9}$ - another coloring


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- A (vertex- or edge-) coloring $c$ of a graph $G$ breaks an automorphism $\varphi$ of $G$ if $\varphi$ does not preserve colors of $c$.
- How many colors we need in a coloring that breaks every non-trivial automorphism of $G$ ?.


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- but may be considered also for edge coloring (Kalinowski, Pilśniak, 2013),
- in both cases the coloring can be proper [proper, vertex -Collins, Trenk 2006]
- Another possibility: endomorphisms instead of automorphisms [W. Imrich, R. Kalinowski, F. Lehner and M. Pilśniak].


## Distinguishing chromatic index

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- Theorem (Kalinowski, Pilśniak, 2013+) Let $G$ be a connected graph of order $n \geq 3$. Then

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except for four graphs of small order: $C_{4}, K_{4}, C_{6}$ or $K_{3,3}$.

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- a contradiction.
$\chi_{D}^{\prime}=3, \mu=4$



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On decomposing regular graphs into locally irregular subgraphs, European Journal of Combinatorics, 49 (2015), 90-104.

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- exact results concerning many classes of graphs,
- Even for trees is not completely solved.


## Irregularity strengh and coloring



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- irregularity strength is minimum $k$ such that there exists an $f$ distinguishing all vertices.


## Irregularity strength: local version

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- ( $G$ connected, $G \neq K_{2}$ )
- $\operatorname{gndi}_{\Sigma} \leq 3$


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## ... 2 conjectures ...

1-2-3 Conjecture gndi $_{\Sigma} \leq 3$

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(L. Addario-Berry, R.E.L. Aldred, K. Dalal, B. A. Reed; 2005)
- gndi $_{\mathcal{M}} \leq 3$ for graphs with large minimum degree (L. Addario-Berry, R.E.L. Aldred, K. Dalal, B. A. Reed; 2005)


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- Such graphs exist for every order $n$.
- We can investigate decompositions of graphs into locally irregular subgraphs.
- Such a decomposition (into $k$ graphs) may be considered as a coloring with $k$ colors such that every color class induces a locally irregular subgraph in $G$.


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- Fortunately, such graphs can be easily characterized:
- paths of odd length,
- cycles of odd length,
- graphs belonging to the family $\mathcal{F}$.

Family $\mathcal{F}$


## Family $\mathcal{F}$; adding even paths

... adding odd paths with a triangle ...


- Theorem. If $G$ is a connected graph, $G \notin \mathcal{F}, G \neq P_{2 p+1} . G \neq C_{2 p+1}$, then it can be decomposed into locally irregular subgraphs.
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- Conjecture Every connected graph $G$, $G \notin \mathcal{F}, G \neq P_{2 p+1} . G \neq C_{2 p+1}$, can be decomposed into 3 locally irregular subgraphs.


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- if an edge $u v \in E$ has color $i$ assigned by a locally irregular edge coloring, then the numbers of edges colored with $i$ incident with $u$ and $v$ must be distinct.


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- For, the graph induced by $i$ is locally irregular.
- So, $u$ and $v$ can be distinguished by multisets of colors


## ... 3 conjectures ...

1-2-3 Conjecture

## gndi $_{\Sigma} \leq 3$

decomposition into 3 locally irregular graphs

## Results

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- Every $d$-regular graph $G$ with $d \geq 10^{7}$ can be decomposed into three locally irregular subgraphs.


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## Complexity

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## Complexity

- Can a graph be decomposed into two locally irregular subgraphs? NP-complete [Julien Bensmail; 2013]
- linear algorithm for determine how many subgraphs we need in the case of trees [O. Baudon, J. Bensmail and E. Sopena, 2013]


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- sums, sets, multisets ... .


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- see for instance: Evelyne FLANDRIN, Hao LI, Antoni MARCZYK, Jean-François SACLÉ, Mariusz WOŹNIAK, A note on neighbor expanded sum distinguishing index, Discussiones Math. - Graph Theory, 2016.


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